AMENDMENTS TO THE CLAIMS:

The following listing of claims supersedes all prior versions and listings of claims in this application:

(Currently Amended) A method of enabling multiple different operating systems to run concurrently on the same computer, <u>said method</u> comprising: selecting a first operating system to have a relatively high priority; selecting at least one second operating system to have a relatively lower priority; providing a common program arranged to switch between said operating systems under predetermined conditions; and

providing modifications to said first and second operating systems to allow them.

providing modifications to said first and second operating systems to allow them to be controlled by said common program;

wherein switching between said operating systems includes invoking the common program by calling an exception vector, and

wherein calling an exception vector to invoke the common program simulates an exception caused by an external event.

2. (Cancelled)

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 (Original) The method of claim 2, comprising allocating exception vectors to trap calls, thereby to enable invocation of the common program using a trap call mechanism.

4-5. (Cancelled)

- (Previously Presented) The method of claim 1, wherein the common program
 preempts the first or second operating system by intercepting exception or interrupt
 vectors.
- 7. (Original) The method of claim 6, further comprising using a exception handler table containing an array of pointers to intercept exceptions, and activating an exception handler program to preempt the first or second operating system.
- (Currently Amended) The method of claim 1, wherein the common program is operable in real mode in physical processor address space.
- (Currently Amended) The method of claim 8, comprising preempting the first or second operating system by the common program, and switching to real mode in

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<u>physical processor address space</u> when preempting the first or second operating system.

- 10. (Currently Amended) The method of claim 8, comprising invoking the common program by the first or second operating system, and switching to real mode <u>in physical processor</u> address space when invoking the common program.
- 11. (Previously Presented) The method of claim 1, comprising enabling hardware interrupts throughout the operation of the second operating system except during the operation of subroutines that save machine state.
- 12. (Previously Presented) The method of claim 1, in which the first operating system is a real time operating system.
- 13. (Previously Presented) The method of claim 1, in which the second operating system is a non-real time, general-purpose operating system.
- 14. (Previously Presented) The method of claim 1, in which the second operating system is Linux, or a version or variant thereof.

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15. (Previously Presented) The method of claim 1, in which the common program is arranged to save, and to restore from a saved version, the processor state

required to switch between the operating systems.

16. (Previously Presented) The method of claim 1, in which processor exceptions for the second operating system are handled in virtual fashion by the

common program.

17. (Previously Presented) The method of claim 1, in which the common program is arranged to intercept some processor exceptions, and to call exception

handling routines of the first operating system to service them.

18. (Original) The method of claim 17, in which the processor exceptions for the

second operating system are notified as virtual exceptions.

19. (Original) The method of claim 18, in which the common program is

arranged to call an exception handling routine of the second operating system

corresponding to a said virtual exception which is pending.

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20. (Previously Presented) The method of claim 1, further comprising providing

each of said operating systems with separate memory spaces in which each can

exclusively operate.

21. (Previously Presented) The method of claim 1, further comprising providing

each of said operating systems with first input and/or output devices of said computer to

which each has exclusive access.

22. (Original) The method of claim 21, in which each operating system

accesses said first input and/or output devices using substantially modified native

routines.

23. (Previously Presented) The method of claim 1, further comprising providing

each of said operating systems with access to second input and/or output devices of

said computer to which each has shared access.

24. (Original) The method of claim 23, in which all operating systems access

said second input and/or output devices using the routines of the first operating system.

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25. (Original) The method of claim 24, in which the common program provides

trap call mechanisms, to control the operation of the second operating system, and/or

event mechanisms to notify the first operating system of status changes in the second

operating system.

26. (Previously Presented) The method of claim 1, further comprising combining

said operating systems and common program into a single code product.

27. (Previously Presented) The method of claim 1, further comprising

embedding said operating systems and common program onto persistent memory on a

computer product.

28. (Previously Presented) A development kit computer program product

comprising code for performing the steps of claim 1.

29. (Currently Amended) A computer program product comprising code

combined according to claim [[36]] 26.

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30. (Currently Amended) A computer system comprising:

a CPU, memory devices and input/output devices, said CPU being arranged to execute having executing thereon computer code comprising[[:1]:

a first operating system having a relatively high priority;

a second operating system having a relatively lower priority; and

a common program arranged to run said operating systems concurrently by switching between said operating systems under predetermined conditions.

wherein switching between said operating systems includes the first or second operating system invoking the common program by calling an exception vector, wherein calling an exception vector to invoke the common program simulates an exception caused by an external event.

- 31. (Currently Amended) A computer system according to claim 30, arranged to run said first and second operating systems concurrently using the method of claim 1 as described above.
- 32. (Currently Amended) The system, product or method of claim 1 in which the computer has a Reduced Instruction Set architecture.